# LANGUAGES RELATED TO THE PROPERTIES OF DISJUNCTIVITY AND CODE 

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#### Abstract

In this paper, we define two new types of words, $d_{1}$-words and $d_{2}$-words and show that the free semigroup $X^{+}$can be represented as a disjoint union of the disjunctive languages $D(1)$, the set of all $d$-primitive words, $N_{d_{1}}$ the set of all $d_{1}$-words and $N_{d_{2}}$ the set of all $d_{2}$-words. Both the languages $N_{d_{1}}$ and $N_{d_{2}}$ are also shown to be disjunctive languages. We show that for any language $L \subseteq X^{+}$, the language $L \cdot Q$, where $Q$ is the set of all primitive words, is not a 2 -code and we give a characterization of that $P \cdot Q^{(i)}$ is a 2-code for some prefix code $P$ and $Q^{(i)}=\left\{f^{i} \mid f \in Q\right\}, i \geq 2$. We proved that each of the disjunctive languages $Q, D(1)^{(i)}, i \geq 1, D(n), n \geq 2, N_{d_{1}}$, and $N_{d_{2}}$ removing a code from it results a dense language. On the other hand, the set $Q^{(i)}, i \geq 2$ removing a code or $Q$ removing a prefix code both results disjunctive languages. Any disjunctive language union a non-dense language is disjunctive. Also any disjunctive language removing a non-dense language is also shown to be disjunctive.


Keywords: Disjunctive language, prefix code, 2 -code, dense language, $d_{1}$-word, $d_{2}$-word

## 1. Introduction

In this paper we assume that $X$ is a finite alphabet containing more than one letter. Let $X^{*}$ be the free monoid generated by $X$. Any element of $X^{*}$ is a word and any subset of $X^{*}$ is a language. For any language $L$, let $|L|$ be the cardinality of $L$. For any two languages $L_{1}, L_{2}$ contained in $X^{*}$, by $L_{1} \subseteq L_{2}$, we mean that $L_{1}$ is a subset of $L_{2}$. We denote it by $L_{1} \subset L_{2}$ when $L_{1} \subseteq L_{2}$ and $L_{1} \neq L_{2}$. Let $X^{+}=X^{*} \backslash\{1\}$, where 1 is the empty word. For any two languages $A$ and $B$, the concatenation of $A$ and $B$ is the set $A B=\{x y \mid x \in A, y \in B\}$ and $A^{+}=A \cup A^{2} \cup A^{3} \cup \ldots$. For a word $w \in X^{*}$, let $\lg (w)$ be the length of the word $w$, which is just the total number

